Single machine problems

- **maximal polynomially solvable:**

$1\|prec;r_i\|C_{max}$	Lawler (1973) [10]
$1\|prec;p_i=p;r_i\|L_{max}$	Simons (1978) [20]
$1\|prec;r_i;pmtn\|L_{max}$	Blazewicz (1976) [6], Baker et al. (1983) [2]
$1\|prec;p_i=p;r_i\|\sum C_i$	Simons (1983) [21]
$1\|prec;pmtn;p_i=p;r_i\|\sum C_i$	Baptiste et al. (2004) [5]
$1\|r_i;pmtn\|\sum C_i$	Baker (1974) [1]
$1\|p_i=p;r_i\|\sum w_iC_i$	Baptiste (2000) [4]
$1\|sp-graph\|\sum w_iC_i$	Lawler (1978) [12]
$1\|r_i;pmtn\|\sum U_i$	Lawler (1990) [13]
$1\|p_i=p;r_i\|\sum w_iU_i$	Baptiste (1999) [3]
$1\|pmtn;p_i=p;r_i\|\sum w_iU_i$	Baptiste (1999) [3]
$1\|p_i=p;r_i\|\sum T_i$	Baptiste (2000) [4]
$1\|pmtn;p_i=p;r_i\|\sum T_i$	Tian et al. (2006) [22]
$1\|p_i=1;r_i\|\sum w_iT_i$	Assignment-problem

- **maximal pseudopolynomially solvable:**

$1\|r_i;pmtn\|\sum w_iU_i$	Lawler (1990) [13]
$1\|\|\sum T_i$	Lawler (1977) [11], Du & Leung (1990) [7]

- **minimal NP-hard:**

*	$1\|r_i\|L_{max}$	Lenstra et al. (1977) [18]
*	$1\|chains;r_i;pmtn\|\sum C_i$	Lenstra (-) [15]
*	$1\|prec\|\sum C_i$	Lawler (1978) [12], Lenstra & Rinnooy Kan (1978) [16]
*	$1\|r_i\|\sum C_i$	Lenstra et al. (1977) [18]
*	$1\|chains;p_i=1;r_i\|\sum w_iC_i$	Lenstra & Rinnooy Kan (1980) [17]
*	$1\|prec;p_i=1\|\sum w_iC_i$	Lawler (1978) [12], Lenstra & Rinnooy Kan (1978) [16]
*	$1\|r_i;pmtn\|\sum w_iC_i$	Labetoulle et al. (1984) [9]
*	$1\|chains;p_i=1\|\sum U_i$	Lenstra & Rinnooy Kan (1980) [17]
	$1\|\|\sum w_iU_i$	Lawler & Moore (1969) [14], Karp (1972) [8]
	$1\|\|\sum T_i$	Lawler (1977) [11], Du & Leung (1990) [7]
*	$1\|chains;p_i=1\|\sum T_i$	Leung & Young (1990) [19]
*	$1\|\|\sum w_iT_i$	Lawler (1977) [11], Lenstra et al. (1977) [18]

- **minimal open:**

$1|pmtn;p_i=p;r_i|\sum w_iC_i$
$1|p_i=p;r_i|\sum w_iT_i$

- **maximal open:**

$1|p_i=p;r_i|\sum w_iT_i$
$1|pmtn;p_i=p;r_i|\sum w_iT_i$

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