

Fig. 5.17. Mapping the two bits 1 of $T_{l}$ to bits 1 of $S$
immediately right of the $(6 n b-9 n)^{\text {th }}$ bit 0 . Since already $3 n$ bits 0 of $S$ have been omitted and right of the $6 n(b-1)^{\text {th }}$ bit 0 of $S^{\prime}$ there are further $3 n$ bits 0 , we know that again at least $3 n$ bits 0 of $S^{\prime}$ are omitted when embedding $T_{l}$ into $S$. Since $T_{l}$ has $15 n^{2}+3 n$ bits $0, S^{\prime}$ has $6 n^{2}+3 n$ bits 0 , and $6 n$ bits 0 of $S$ are not used when embedding $T_{l}$ into $S$, we know that the remaining suffix of $T_{l}$ must be embedded into $S^{\prime \prime}$. Thus $S^{\prime \prime}$ contains at least $6 n+9 n^{2}$ bits 0 . Denote the exact number of bits 0 in $S^{\prime \prime}$ by $p+9 n^{2}$ with some number $p \geq 6 n$. Since each of the strings $A_{i}$ and $S^{\prime}$ contain the same number $6 n^{2}+3 n$ of bits 0 we know that all of the strings $B_{j}$ must be embedded into $S^{\prime \prime}$. By Lemma 5.20 we know that $S^{\prime \prime}$ must contain at least $\left(9 n^{2}+1\right) /(p+1)$ many bits 1 . Thus $S^{\prime \prime}$ has length at least $p+9 n^{2}+\left(9 n^{2}+1\right) /(p+1)$. Since $K=15 n^{2}+10 n+k$ was the length of $S$ and $S^{\prime}$ contained $6 n^{2}+3 n$ bits 0 and at least $n$ bits 1, we conversely know that the length of $S^{\prime \prime}$ can be at most $9 n^{2}+6 n+k$. This results in the following inequality which we lead to a contradiction (using $k \leq n$ ) by a sequence of transformation steps:

$$
\begin{align*}
& p+9 n^{2}+\frac{9 n^{2}+1}{p+1} \leq 9 n^{2}+6 n+k \leq 9 n^{2}+7 n \\
\Leftrightarrow & p+\frac{9 n^{2}+1}{p+1} \leq 7 n \tag{i}
\end{align*}
$$

Now we show

$$
\begin{equation*}
6 n+\frac{9 n^{2}+1}{6 n+1} \leq p+\frac{9 n^{2}+1}{p+1} \tag{ii}
\end{equation*}
$$

by the following equivalence transformations.

